ARCHIMEDEAN LOCAL HEIGHT DIFFERENCES ON ELLIPTIC CURVES

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Abstract. To compute generators for the Mordell-Weil group of an elliptic curve over a number field, one needs to bound the difference between the naive and the canonical height from above. We give an elementary and fast method to compute an upper bound for the local contribution to this difference at an archimedean place, which typically gives better results than previous algorithms.

1. Introduction

Let $E$ be an elliptic curve defined over a number field $K$. By the Mordell-Weil theorem the $K$-rational points on $E$ form a finitely generated group

$$E(K) \cong \mathbb{Z}^r \times \text{Tor}(E(K));$$

where $r \geq 0$ is the rank of $E/K$ and $\text{Tor}(E(K))$ is the (finite) torsion subgroup of $E(K)$. One of the fundamental computational problems in the study of the arithmetic of elliptic curves is to compute generators for $E(K)$. Applications of this include, for instance, the numerical verification of the full conjecture of Birch and Swinnerton-Dyer in examples, as well as the computation of $S$-integral points on $E$, e.g. using the recent approach of von Känel and Matschke [vKM16].

Generators of $\text{Tor}(E(K))$ are typically easy to find. No effective method for the computation of $r$ is known, but there are still several methods which often succeed in practice. Suppose that we know $r$ and $Q_1, \ldots, Q_r \in E(K)$ whose classes generate a finite index subgroup of $E(K)/\text{Tor}(E(K))$. The final step is then to deduce generators of $E(K)$ from this. This is done by saturating the lattice generated by $Q_1, \ldots, Q_r$ inside the Euclidean vector space $(E(K) \otimes \mathbb{R}, \hat{h})$, where $\hat{h}$ is the canonical height. The most widely used saturation algorithm is due to Siksek [Sik95] and requires, in particular, an algorithm to enumerate points on $E(K)$ of canonical height bounded by a fixed real number $B$ (this set is finite by the Northcott property).

In practice, this is done by first computing an upper bound $\beta$ for the difference between $\hat{h}$ and the naive height $h : E(K) \to \mathbb{R}$; the points with canonical height bounded by $B$ are then contained in

$$\{ P \in E(K) : h(P) \leq B + \beta \},$$

which can be enumerated for reasonably small $B + \beta$. Note that the heights we consider are logarithmic, so that $\beta$ shows up exponentially in the size of the search space. It is therefore of great practical importance to make $\beta$ as small as possible. At the same time, it is desirable to keep the computation of $\beta$ reasonably fast.
The standard approach for bounding the difference \( h - \hat{h} \) is to write it as a sum of local terms, one for each place of \( K \), and to bound the local contributions individually, see [CPS06] or Section 3 below. For non-archimedean places, optimal bounds are given in [CPS06]. Our main contribution is Theorem 4.2, which provides an elementary method for bounding the local contribution at an archimedean place. This method is extremely fast in practice, and yields better results than other existing approaches in many examples. The approach is analogous to an algorithm due to Stoll [Sto99], with modifications by Stoll and the first-named author [MS16b] for Jacobians of genus 2 curves and to Stoll [Sto17] for Jacobians of hyperelliptic genus 3 curves. In the case of elliptic curves, the validity of our formulas can be established using essentially only linear algebra.

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2. Action of the two-torsion subgroup

In this section, we let \( K \) be an algebraically closed field of characteristic zero and let \( E/K \) be an elliptic curve, given by a Weierstrass equation

\[
y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6,
\]

with point at infinity \( O \). We denote by \( b_2, \ldots, b_8 \) the usual \( b \)-invariants of \( E \). Let \( \kappa : E \to \mathbb{P}^1 \) be the \( x \)-coordinate map with respect to the given equation (2.1), extended to all of \( E \) by setting \( \kappa(O) = (1 : 0) \). Given a representative \((x_1, x_2)\) for \( \kappa(P) \), we have

\[
\kappa(2P) = \delta(x_1, x_2),
\]

where \( \delta = (\delta_1, \delta_2) \), and

\[
\begin{align*}
\delta_1(x_1, x_2) &= x_1^4 - b_4x_1^2x_2^2 - 2b_6x_1x_2^3 - b_8x_2^4, \\
\delta_2(x_1, x_2) &= 4x_1^3x_2 + b_2x_1^2x_2^2 + 2b_4x_1x_2^3 + b_6x_2^4.
\end{align*}
\]

The purpose of the present section is to prove an explicit version of the following result.

**Proposition 2.1.** There are quadratic forms \( y_1, y_2, y_3 \in K[x_1, x_2] \) and constants \( a_{ij}, b_{jk} \in K \), depending only on \( E \), such that for \( i = 1, 2 \) and \( j = 1, 2, 3 \) we have

\[
x_i^2 = \sum_{j=1}^3 a_{ij}y_j(x_1, x_2) \quad \text{and} \quad y_j(x_1, x_2)^2 = \sum_{k=1}^2 b_{jk}\delta_k(x_1, x_2)
\]

in \( K[x_1, x_2] \).

The constants \( a_{ij} \) and \( b_{jk} \) are given in (2.2) and (2.3), respectively.

For \( T \in E[2] \) let \(+_T : E \to E\) be translation by \( T \). Since \(- (T + P) = T - P\), the map \(+_T\) descends to a map on \( \mathbb{P}^1 \). In fact there is a linear transformation \( m_T \) on \( \mathbb{P}^1 \) such that \( \kappa \circ +_T = m_T \circ \kappa \). A simple calculation shows that \( m_T \) is represented any non-trivial scalar multiple of the matrix

\[
M_T := \begin{cases} 
E_2 & T = O \\
(x(T) & f'(x(T)) - x(T)^2) \\
1 & -x(T)
\end{cases} , \quad T \neq O
\]
where \( f = x^3 + \frac{b_2}{4}x^2 + \frac{b_4}{2}x + \frac{b_6}{4} \).

For the proof of Proposition 2.1, we analyze the action of \( E \) on the space of homogeneous polynomials in two variables of degree 2 and 4, respectively. We first lift the transformation matrices \( M \) to a subgroup \( G_E \) of \( \text{SL}_2(K) \) such that \( E[2] \cong G_E/\{\pm E_2\} \). Let \( e_2 : E[2] \times E[2] \to \mu_2 \) denote the Weil pairing; then \( e_2(T, T') = \varepsilon(T)x(T')\varepsilon(T + T') \), where \( \varepsilon(O) := 1 \) and \( \varepsilon(T) := -1 \) for \( T \in E[2] \setminus \{O\} \).

**Lemma 2.2.** Let \( T, T' \in E[2] \). Then we have

\[
M_{T'}^*M_T = \begin{cases} 
\varepsilon(T) \det(M_T)E_2, & T = T' \\
\varepsilon_2(T, T')M_T M_{T'}, & T \neq T'.
\end{cases}
\]

**Proof.** The assertion is trivial when \( T = O \) or \( T' = O \). Suppose that \( T \neq O \) and \( T' \neq O \). It is easy to compute

\[
M_{T'}^*M_T = \left(\begin{array}{cc}
(x(T)x(T') + 2x(T')^2 + \frac{b_4}{2}x(T') + \frac{b_6}{4} & \varepsilon(T - x(T'))(2x(T)x(T') - \frac{b_4}{2}) \\
x(T) - x(T') & 2x(T)^2 + \frac{b_4}{2}x(T) + \frac{b_6}{4}
\end{array}\right) .
\]

In particular, \( M_T^2 = -\det(M_T)E_2 \). If \( T \) and \( T' \) are distinct, the group law on \( E \) shows

\[
M_{T'}^*M_T = (x(T) - x(T'))^2 \begin{pmatrix}
x(T + T') & 2x(T)x(T') - \frac{b_4}{2} \\
1 & -x(T + T')
\end{pmatrix}
= (x(T) - x(T'))^2 \begin{pmatrix}
x(T + T') & f'(x(T + T')) - x(T + T')^2 \\
1 & -x(T + T')
\end{pmatrix} = -M_T M_{T'},
\]

which proves the result. \( \Box \)

Lemma 2.2 shows that the classes of the matrices \( M_T \) form a subgroup of \( \text{PSL}_2(K) \). We now lift this to a double cover which is a subgroup of \( \text{SL}_2(K) \).

**Lemma 2.3.** For \( T \in E[2] \) let \( \gamma_T \in K^\times \) such that \( \gamma_T^2 = \det(M_T)^{-1} \) and let \( \tilde{M}_T := \gamma_T M_T \). Then

\[
G_E := \{\pm \tilde{M}_T \mid T \in E[2]\}
\]

is a subgroup of \( \text{SL}_2(K) \). Moreover, \( G_E \) is isomorphic to the quaternion group \( Q_8 \), and \( E[2] \cong G_E/\{\pm E_2\} \).

**Proof.** Obviously \( G_E \subseteq \text{SL}_2(K) \) and \( G_E \) does not depend on the choice of \( \gamma_T \). By Lemma 2.2, we have \( M_T^{-1} \in G_E \) for \( M \in G_E \). Let \( T_1, T_2, T_3 \in E[2] \) be nontrivial and pairwise distinct. Since

\[
\kappa \circ T_3 = \kappa \circ T_1 \circ T_2 = m_{T_1} \circ \kappa \circ T_2 = m_{T_1} \circ m_{T_2} \circ \kappa,
\]

we have

\[
M_{T_1} M_{T_2} = \gamma M_{T_3}
\]

for a unit \( \gamma \in K^\times \). From \( \det(\tilde{M}_{T_1} \tilde{M}_{T_2}) = \det(\gamma_{T_1} \gamma_{T_2} \gamma M_{T_3}) = 1 \) we deduce that \( \gamma_{T_3} \) is equal to \( \gamma_{T_1} \gamma_{T_2} \gamma \) up to sign, so \( \tilde{M}_{T_1} \tilde{M}_{T_2} G_E \), and hence \( G_E \) is indeed a subgroup of \( \text{SL}_2(K) \).

The remaining statements are clear. \( \Box \)
In other words, we have a commutative diagram
\[
\begin{array}{cccccc}
0 & \longrightarrow & \mu_2 & \longrightarrow & \mathbb{G}_E & \longrightarrow & \mathbb{E}[2] & \longrightarrow & 0 \\
\downarrow & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\
0 & \longrightarrow & \mu_2 & \longrightarrow & \text{SL}_2 & \longrightarrow & \text{PSL}_2 & \longrightarrow & 0 \\
\end{array}
\]

Proof of Proposition 2.1. Let \( \rho \) denote the standard representation of \( \mathbb{G}_E \) on the vector space \( V \) of \( K \)-linear forms in \( x_1, x_2 \). Then the symmetric square \( \rho^2 \) factors through \( \mathbb{E}[2] \). Hence we can view \( \rho^2 \) as a representation of \( \mathbb{E}[2] \) on \( \text{Sym}^2(V) \), and we have

\[
\rho^2 = \bigoplus_{T \in \mathbb{E}[2] \setminus \{O\}} e_2(\cdot, T).
\]

It is easy to check that for each nontrivial 2-torsion point \( T \) the polynomial

\[
y_T := x_1^2 - 2x(T)x_1x_2 - (f'(x(T)))x_2^2 \in \text{Sym}^2(V)
\]

is an eigenform of \( \rho^2 \). Fix any ordering of the non-trivial 2-torsion points and call them \( T_1, T_2, T_3 \); let \( y_j := Y_{T_j} \). Since \( Y := (y_1, y_2, y_3) \) is linearly independent, \( Y \) forms a basis for \( \text{Sym}^2(V) \). We find that the coefficients of \( x_1^2 \) and \( x_2^2 \) with respect to \( Y \) are given by

\[
I_Y(x_1^2) = \tau^{-1} \begin{pmatrix} f'(x(T_1)) - x(T_1)^2 \\ f'(x(T_2)) - x(T_2)^2 \\ f'(x(T_3)) - x(T_3)^2 \end{pmatrix} \times \begin{pmatrix} x(T_1) \\ x(T_2) \\ x(T_3) \end{pmatrix}
\]

and

\[
I_Y(x_2^2) = \tau^{-1} \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \times \begin{pmatrix} x(T_1) \\ x(T_2) \\ x(T_3) \end{pmatrix},
\]

where \( \tau := \sum_i f'(x(\sigma^i(T_1))) - x(\sigma^i(T_1))^2 \left( x(\sigma^i(T_2)) - x(\sigma^i(T_3)) \right) \neq 0 \) for the cycle \( \sigma = (T_1 T_2 T_3) \). In other words, we have \( x_i^2 = \sum_{j=1}^3 a_{ij} y_j(x_1, x_2) \) for \( i = 1, 2 \), where

\[
\begin{pmatrix}
a_{11} & a_{12} & a_{13} \\
a_{21} & a_{22} & a_{23}
\end{pmatrix}
\]

\[
= \begin{pmatrix} 2x(T_2)x(T_3) - b_4 \\ 2x(T_1)x(T_3) - b_4 \\ 2x(T_1)x(T_2) - b_4 \\ \frac{2x(T_2)x(T_3)}{2(x(T_1) - x(T_2))} - x(T_1) \\ \frac{2x(T_1)x(T_3)}{2(x(T_2) - x(T_3))} - x(T_2) \\ \frac{2x(T_1)x(T_2)}{2(x(T_3) - x(T_1))} - x(T_3) \end{pmatrix}.
\]

As for \( \rho^4 \), we have that \( \rho^4 \) factors through \( \mathbb{E}^2 \). Since projectively \( \delta(M_T(x_1, x_2)) = \delta(x_1, x_2) \), the polynomials \( \delta_1, \delta_2 \) are \( \mathbb{E}^2 \)-invariant under the fourth symmetric power \( \rho^4 \). Moreover, they are linearly independent. As the space of \( \mathbb{E}^2 \)-invariant quartic polynomials is 2-dimensional, it is spanned by \( \delta_1 \) and \( \delta_2 \). Computing the squares \( y_j^2 \), we find that \( y_j(x_1, x_2)^2 = \sum_{k=1}^2 b_{jk} \delta_k(x_1, x_2) \), where

\[
\begin{pmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \\ b_{31} & b_{32} \end{pmatrix} = \begin{pmatrix} 1 & -x(T_1) \\ 1 & -x(T_2) \\ 1 & -x(T_3) \end{pmatrix}.
\]

This completes the proof of Proposition 2.1. \( \square \)
3. Global height differences

Let $K$ be a number field. We define $M_K$ to be the set of places of $K$, where we normalize the absolute value $|\cdot|_v$ associated to $v \in M_K$ by requiring that it extends the usual absolute value on $\mathbb{Q}$ when $v$ is an infinite place and by setting $|p|_v = p^{-1}$ when $v$ is a finite place above a prime number $p$. For $v \in M_K$, we set $n_v = [K_v : \mathbb{Q}_w]$ where $w$ is the place of $\mathbb{Q}$ below $v$. Then the product formula $\prod_{v \in M_K} |x|_v^{|\nu_v|} = 1$ holds for all $x \in K^\times$.

Consider an elliptic curve $E/K$, given by an integral Weierstrass equation (2.1). We define the naive height of $P \in E(K) \setminus \{O\}$ by

$$h(P) = h(\kappa(P)) = \frac{1}{[K : \mathbb{Q}]} \sum_{v \in M_K} n_v \log \max\{|x_1|_v, |x_2|_v\},$$

where $(x_1, x_2) \in K^2$ represents $\kappa(P)$. The canonical height of $P$ is defined as the limit

$$\hat{h}(P) = \lim_{n \to \infty} 4^{-n} h(2^n P).$$

In this work, we are not really interested in the canonical height itself, but rather in upper bounds on the difference $h - \hat{h}$. As in [CPS06] and [MS16a], we decompose the difference into a finite sum of local terms

$$h(P) - \hat{h}(P) = \frac{1}{[K : \mathbb{Q}]} \sum_{v \in M_K} n_v \Psi_v(P),$$

where the functions $\Psi_v : E(K_v) \to \mathbb{R}$ are continuous and bounded. It is then clear that it suffices to compute upper bounds on all $\Psi_v$ to deduce an upper bound on the difference $h - \hat{h}$. Recall that we may write

$$\Psi_v(P) = -\sum_{n=0}^{\infty} 4^{-n-1} \log \Phi_v(2^n P)$$

for $P \in E(K_v)$, where $\Phi : E(K_v) \to \mathbb{R}$ is the continuous bounded function defined by

$$\Phi_v(P) := \max\{|\delta_1(x_1, x_2)|_v, |\delta_2(x_1, x_2)|_v\}$$

and $(x_1, x_2) \in K_v^2$ represents $\kappa(P)$. See [CPS06] and [MS16a] for details.

4. Archimedean local height differences

In this section we show how to bound the local contribution $\Psi_v$ to the height difference, where $v$ is an archimedean place of a number field. We will drop $v$ from the notation for simplicity and assume that $K_v = \mathbb{C}$, unless stated otherwise. So consider an elliptic curve $E/\mathbb{C}$, given by a Weierstrass equation (2.1). Note that Proposition 2.1 lets us bound $|a_i|^4$ ($i = 1, 2$) in terms of $|\delta_j(x_1, x_2)|$, $j = 1, 2$. From this we easily get an upper bound for $\Phi$ using the triangle inequality. Via the geometric series we deduce:

**Corollary 4.1.** We have

$$\max_{P \in E(\mathbb{C})} \{\Psi(P)\} \leq \frac{4}{3} \left( \sqrt[3]{\sum_{j=1}^{3} |a_{ij}| \sqrt{|b_{j1}| + |b_{j2}|}} \right)_{i=1,2}.$$
where the constants $a_{ij}, b_{jk} \in \mathbb{C}$ are as in Proposition 2.1.

This idea was first used by Stoll [Sto99, Sto17] to bound the height difference for Jacobians of genus 2 curves and hyperelliptic genus 3 curves, respectively. We will follow his approach closely; in fact, the elliptic case is much simpler. Furthermore, we will iterate the bound for $\Phi$ to get a better bound for $\Psi$ than the one obtained from the geometric series; this was used by Stoll and the first-named author for genus 2 [MS16b], and by Stoll [Sto17] for genus 3.

For the iteration we define the function

$$
\varphi : \mathbb{R}_+^2 \to \mathbb{R}_+^2, (d_1, d_2) \mapsto \left( \sum_{j=1}^3 \left| a_{ij} \sqrt{|b_{j1}| d_1 + |b_{j2}| d_2} \right|, i = 1, 2 \right)
$$

and we set

$$
c_N := \frac{4^N}{4^N - 1} \log(\|\varphi \circ N(1, 1)\|)
$$

for $N \geq 1$, where $\| \cdot \|$ denotes the supremum norm. Hence $c_1$ is precisely the upper bound from Corollary 4.1. Our algorithm for bounding $\Psi$ is based on the following result, whose statement and proof follow [MS16b, Lemma 16.1].

**Theorem 4.2.** The sequence $(c_N)_{N \geq 1}$ is monotonically decreasing and we have

$$
\max_{P \in E(\mathbb{C})} \{\Psi(P)\} \leq c_N
$$

for every $N \geq 1$.

**Proof.** To verify that the upper bound holds, let $\alpha \in \mathbb{C}^2$. A simple induction shows that for $N \geq 1$ we have

$$
|\alpha_i| \leq \varphi \circ N \left( |\delta \circ N(\alpha)_1|, |\delta \circ N(\alpha)_2| \right)_i,
$$

and since

$$
|\delta \circ N(\alpha)_i| \leq \varphi(1, 1)_i \|\delta \circ (N+1)(\alpha)\|^{\frac{1}{4}},
$$

we find that

$$
|\alpha_i| \leq \varphi \circ N \left( \|\delta \circ (N+1)(\alpha)\|^{\frac{1}{4}} \varphi(1, 1) \right)_i.
$$

Shifting $N$ by 1 and using that $\varphi$ is homogeneous of degree 1/4, it follows that

(4.1) $$
|\alpha_i| \leq \|\delta \circ N(\alpha)\|^{\frac{1}{4N}} \varphi \circ N(1, 1)_i.
$$

We now apply (4.1) to $\alpha = \delta \circ N^n(x_1, x_2)$, where $n \geq 1$ and $x \in \mathbb{C}^2$ represents $\kappa(P)$ for $P \in E(\mathbb{C})$, to obtain

$$
\|\delta \circ N^n(x_1, x_2)\| \leq \|\delta \circ N(n+1)(x)\|^{\frac{1}{4N}} \|\varphi \circ N(1, 1)\|.
$$

Upon noting that

$$
\Psi(P) = \sum_{n=0}^{\infty} 4^{-Nn} \log \left( \frac{\|\delta \circ N^n(x_1, x_2)\|}{\|\delta \circ N(n+1)(x_1, x_2)\|^{\frac{1}{4N}}} \right)
$$

the result follows.

To show that $c_N$ is monotonically decreasing, consider the function

$$
\psi : \mathbb{R}^2 \to \mathbb{R}^2, \alpha \mapsto \left( \log(\varphi(\exp(\alpha_1), \exp(\alpha_2)))_i \right)_{i=1,2}.
$$
Note that the Jacobi matrix of $\psi$ has positive entries and that its rows sum to $1/4$, because $\varphi_1$ and $\varphi_2$ are homogeneous of degree $1/4$. It follows that

$$\|\psi(\alpha) - \psi(\beta)\| \leq \frac{1}{4}\|\alpha - \beta\|$$

for all $\alpha, \beta \in \mathbb{C}^2$.

For $N \geq 1$ we have $c_N = \frac{4}{4^N - 1}\|\psi^N(0,0)\|$. In particular, (4.2) implies

$$\|\psi^2(0,0) - \psi(0,0)\| \leq \frac{1}{4}\|\psi(0,0)\|,$$

whence $c_2 \leq c_1$. We now proceed by induction on $N$; so let $N \geq 2$ such that $c_N \leq c_{N-1}$. Let $b_N := \|\psi^N(0,0)\|$. If $b_{N+1} \leq b_N$, then we’re done, so we may assume that $b_{N+1} > b_N$. Applying (4.2) to $\alpha = \psi^N(0,0)$ and $\beta = \psi^{(N-1)}(0,0)$, we find that

$$b_{N+1} \leq \min \left\{ \frac{5}{4}b_N - \frac{1}{4}b_{N-1}, \frac{3}{4}b_N + \frac{1}{4}b_{N-1} \right\},$$

according to whether $b_N \geq b_{N-1}$ or not. First assume that $b_N \geq b_{N-1}$, so that

$$b_{N+1} \leq \frac{5}{4}b_N - \frac{1}{4}b_{N-1}.$$

From $c_N \leq c_{N-1}$ we get

$$-\frac{1}{4}b_{N-1} \leq -\frac{4^{N-1} - 1}{4^{N-1} - 4}b_N,$$

which implies

$$b_{N+1} \leq \left( \frac{5}{4} - \frac{4^{N-1} - 1}{4^{N-1} - 4} \right) b_N = \frac{4^{N+1} - 2}{4^{N+1} - 4}b_N < \frac{4^{N+1} - 1}{4^{N+1} - 4}b_N,$$

and hence $c_{N+1} < c_N$. The case $b_N < b_{N-1}$ is similar. \qed

In particular, $(c_N)_N$ and $(b_N)_N$ both converge to the same limit, and this limit is an upper bound for $\Psi$. In practice, the sequence converges quickly, and a few iterations suffice. This gives us a very simple method to bound $\Psi$ from above.

**Remark 4.3.** Suppose that $v$ is a real place and that $E(\mathbb{R})$ has only one component. Then $b_{22}$ and $b_{32}$ are non-real, but all $P \in E(K_v)$ have real coordinates, so we have

$$|y_j(x_1, x_2)^2| = \sum_{k=1}^2 b_{jk}^j\delta_k(x_1, x_2)$$

$$\leq \max \{|b_{j1}\delta_1(x_1, x_2) + b_{j2}\delta_2(x_1, x_2)|, |b_{j1}\delta_1(x_1, x_2) - b_{j2}\delta_2(x_1, x_2)|\}$$

for $j \in \{2,3\}$ and $x \in \mathbb{R}^2$ representing $\kappa(P)$. Modifying the definition of the function $\varphi$ accordingly, we often get a better bound in practice.
5. Alternative algorithms

In this section we briefly discuss other approaches to bounding $\Psi_v$ from above for an archimedean place $v$. The approach of Cremona-Prickett-Siksek [CPS06] is to find the largest value $\gamma$ of $\Phi_v$; then $\gamma/3$ is an upper bound for $\Psi_v$. For real places this translates into a simple algorithm which is trivial to implement. For complex places, they give two approaches: one based on Gröbner bases and another one based on refining an initial crude bound via repeated quadrisection. The latter is faster and yields better bounds in practice than the former. The method of [CPS06] is implemented in Magma and as part of Cremona’s mwrank (which is also contained in Sage). A variation of this approach was presented by Uchida [Uch08]; he computes the largest value of an analogue of $\Phi_v$, but with duplication replaced by multiplication by $m$ for $m > 2$.

An alternative approach is to use that for $K_v = \mathbb{C}$, which we may assume without loss of generality, $\Psi_v$ can be expressed in terms of the Weierstrass $\wp$-function and an archimedean canonical local height function, which in turn is closely related to the Weierstrass $\sigma$-function. This was used by Silverman [Sil90] to provide an easily computed upper bound for $\Psi_v$ in terms of the values of the $j$-invariant and the discriminant of $E$; according to [CPS06], this bound is usually larger than the one due to Cremona-Prickett-Siksek, at least for real embeddings. In a spirit similar to the repeated quadrisection method in [CPS06], Bruin [Bru13] uses a recursive approach (starting from a fundamental domain of the period lattice of $E/\mathbb{C}$) to approximate the maximal value taken by $\Psi_v$ on $E(\mathbb{C})$ to any desired precision. Bruin’s algorithm therefore gives nearly optimal bounds for complex embeddings, whereas for real embeddings the bound computed using the algorithm of Cremona-Prickett-Siksek is often smaller. A Pari/GP implementation of Bruin’s method can be found at https://www.math.leidenuniv.nl/~pbruin/hdiff.gp (note that this uses a different normalization from ours; the difference is $\log |\Delta|_v/6$, where $\Delta$ is the discriminant of the given Weierstrass model). While this method is reasonably fast for curves with small coefficients, it can be slow even for medium-sized coefficients. For instance, it took about 18 minutes to compute an upper bound for the curve with Cremona label 11a2, which has minimal Weierstrass equation

$$y^2 + y = x^3 - x^2 - 7820x - 263580.$$  

So while this approach leads to superior bounds, it is somewhat less useful in practice, because the need for computing a very sharp upper bound mostly arises for curves whose coefficients are relatively large.

6. Experiments and comparison

We implemented an algorithm based on Theorem 4.2 and Remark 4.3 to compute an upper bound for $\Psi_v$ for an archimedean place $v$ in Magma [BCP97]. The code is available at https://github.com/steffenmueller/arch-ht-diff. We experimentally compared our code with the Magma-implementation of the algorithm of [CPS06], using a single core on an Intel Xeon(R) CPU E3-1275 V2 3.50GHz processor. Note that the latter sometimes shows that the upper bound is exactly 0 (which is attained by $P = O$), whereas our code always returns a positive real number. We compared all curves of conductor at most 35,000 in Cremona’s database of elliptic curves over the rationals. Here and in
the following $\beta$ is the upper bound returned by our code and $\beta_{CPS}$ is the upper bound returned by the Magma implementation of the algorithm from [CPS06]. We also list the average value of $\beta$ and $\beta_{CPS}$ (including the cases where the latter is 0).

<table>
<thead>
<tr>
<th>max. conductor</th>
<th>$\beta_{CPS} = 0$</th>
<th>$\beta &gt; \beta_{CPS}$</th>
<th>$\beta &lt; \beta_{CPS}$</th>
<th>avg. $\beta_{CPS}$</th>
<th>avg. $\beta$</th>
</tr>
</thead>
<tbody>
<tr>
<td>10,000</td>
<td>33.5%</td>
<td>38.8%</td>
<td>27.8%</td>
<td>0.947</td>
<td>0.992</td>
</tr>
<tr>
<td>20,000</td>
<td>33.7%</td>
<td>37.9%</td>
<td>28.3%</td>
<td>0.979</td>
<td>1.007</td>
</tr>
<tr>
<td>35,000</td>
<td>33.8%</td>
<td>37.5%</td>
<td>28.8%</td>
<td>1.001</td>
<td>1.007</td>
</tr>
</tbody>
</table>

We found similar results for databases of ‘small’ elliptic curves over real quadratic fields. Perhaps surprisingly, the picture is quite different for ‘random’ curves, and it would be interesting to investigate why this is the case. The following table contains the respective results for $10^5$ randomly chosen elliptic curves over $\mathbb{Q}$ with Weierstrass coefficients $a_1, \ldots, a_6$ bounded in absolute value by $B \in \{10^2, 10^3, 10^4\}$.

<table>
<thead>
<tr>
<th>$B$</th>
<th>$\beta_{CPS} = 0$</th>
<th>$\beta &gt; \beta_{CPS}$</th>
<th>$\beta &lt; \beta_{CPS}$</th>
<th>avg. $\beta_{CPS}$</th>
<th>avg. $\beta$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$10^2$</td>
<td>46.3%</td>
<td>3.3%</td>
<td>50.4%</td>
<td>0.145</td>
<td>0.045</td>
</tr>
<tr>
<td>$10^3$</td>
<td>48.4%</td>
<td>1.0%</td>
<td>50.7%</td>
<td>0.146</td>
<td>0.011</td>
</tr>
<tr>
<td>$10^4$</td>
<td>49.2%</td>
<td>0.3%</td>
<td>50.5%</td>
<td>0.147</td>
<td>0.002</td>
</tr>
</tbody>
</table>

In the above tables, the average time it took to compute the bounds was very short (less than 0.002 seconds on average) for both algorithms. A comparison over $\mathbb{Q}(\sqrt{5})$ with the same parameters resulted in the following:

<table>
<thead>
<tr>
<th>$B$</th>
<th>$\beta_{CPS} = 0$</th>
<th>$\beta &gt; \beta_{CPS}$</th>
<th>$\beta &lt; \beta_{CPS}$</th>
<th>avg. $\beta_{CPS}$</th>
<th>avg. $\beta$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$10^2$</td>
<td>21.4%</td>
<td>6.8%</td>
<td>71.9%</td>
<td>0.148</td>
<td>0.039</td>
</tr>
<tr>
<td>$10^3$</td>
<td>23.5%</td>
<td>1.9%</td>
<td>74.6%</td>
<td>0.148</td>
<td>0.010</td>
</tr>
<tr>
<td>$10^4$</td>
<td>24.6%</td>
<td>0.4%</td>
<td>75.0%</td>
<td>0.150</td>
<td>0.002</td>
</tr>
</tbody>
</table>

So it seems that for large coefficients, our algorithm yields better results most of the time, unless $\beta_{CPS} = 0$. We also see that, on average, our bound is much smaller. Here is a particularly striking example.

**Example 6.1.** Let $E/\mathbb{Q}$ be given by

$$y^2 + xy + y = x^3 - x^2 + 31368015812338065133318565292206590792820353345x$$
$$+ 30203880269856608735643188429543498624522041683874493555186062568159847$$

This example was found by Elkies in 2009 and currently holds the record for the elliptic curve of largest known rank ($r = 19$) which is provably correct, independently of any conjectures. In this case $\beta_{CPS} = 18.018$, whereas $\beta = 0.147$.

We also compared the two implementations for a few thousand curves with coefficient sizes as above, but over some imaginary quadratic fields. Here we found that our bound was better in all examples. Moreover, it also took less time to compute in all cases (on average 0.003 seconds compared to 1.2 seconds).

In practice, the methods of this paper, of Cremona-Prickett-Siksek and of Bruin should be combined. For a real embedding, one should first compute an upper bound using Cremona-Prickett-Siksek. If this is non-zero, one should then apply our algorithm, and use whichever bound is smaller. For a complex embedding, our algorithm appears to be a good first choice. If the resulting bound seems too large for saturation, and if the
coefficients of the curve are of reasonable size, one can then compute a bound using the algorithm of Bruin. This is basically optimal for complex embeddings, and sometimes beats the other bounds for real embeddings as well, but, as discussed above, it typically takes much longer to compute.

References


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