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Nested grids ILU-decomposition (NGILU)

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Abstract

A preconditioning technique is described which shows, in many cases, grid-independent convergence. This technique only requires an ordering of the unknowns based on the different levels of multigrid, and an incomplete LU-decomposition based on a drop tolerance. The method is demonstrated on a variety of well-known elliptic test problems including strongly varying coefficients, advective terms and grid refinement.

Keywords: Finite differences; Preconditioning; Multilevel methods; Conjugate gradients; Incomplete decomposition

AMS classification: 65F10; 65F50; 65N22; 65N55

1. Introduction

Solution of large sparse systems of linear equations continues to be a major research area with widespread application. In this paper, we consider linear systems of the form \(Ax = b\), where \(A\) is a large \(N \times N\) matrix and \(b\) a given vector. Elimination methods have as drawback that, in general, the number of nonzero entries strongly increases during the elimination process. Therefore, these methods are often too costly in terms of computer storage and CPU-time. As a consequence, sparse systems of linear equations are often solved by an iterative method, for example by some conjugate gradient-like method. In general, the speed of convergence of these methods strongly depends on the eigenvalue distribution of \(A\). For example, for the conjugate gradient method (CG) it can be proven that it is very important that the spectral condition number is small, and that the extreme eigenvalues are well separated [17]. Therefore, a CG-like method is often applied to the preconditioned system \((LU)^{-1}Ax = (LU)^{-1}b\) instead of the original system \(Ax = b\). Herein \(L\) and \(U\) are
sparse factors obtained by some incomplete LU-decomposition of $A$. The matrix $L$ is a lower-triangular matrix, and $U$ is upper-triangular.

Meijerink and van der Vorst [10] have shown existence and uniqueness of the incomplete LU-decomposition for an important class of problems and for an arbitrary choice of the sparsity pattern of $L + U$. For symmetric systems, they showed that CG gives excellent results when combined with an incomplete Choleski-decomposition as preconditioner (ICCG). Gustafsson [8] has shown that, in many cases, the performance of an incomplete LU-decomposition can be improved by lumping all elements that are neglected during the decomposition on the main diagonal. For several problems, the CPU-time for the iteration combined with the resulting Modified ILU-decomposition (MILU) is $O(N^{5/4})$ in two dimensions and $O(N^{7/6})$ for 3D-problems, where $N$ is the total number of unknowns (see also [2]).

Multigrid methods perform even better, and for a large class of problems they have an optimal order of convergence: the amount of work and storage is proportional to $N$. However, due to the required proper smoothers and the restriction and prolongation operators at each level, the implementation of multigrid techniques for practical problems is much more complicated than that of some CG-like method combined with a preconditioner. We present a method which combines the best properties of both, e.g. an incomplete LU-decomposition such that the preconditioned system can be solved with the optimal computational complexity $O(N)$ by some CG-like method. The basic idea is the same as in multigrid methods, which use coarser grids in order to remove the low-frequency errors effectively. Our preconditioning technique uses a partitioning of the unknowns based on the sequence of grids in multigrid.

In the special case when $A$ comes from a standard discretisation of a stationary convection–diffusion equation on a uniform, rectangular grid, the partitioning of the unknowns is the same as in the so-called RRB-method described in [4], the nested recursive two-level decomposition method described in [1], and the method described in [11]. The lumping strategy in these methods is based on an approximation of nine-point stencils with five-point stencils. In [4] it is shown that, when $A$ comes about by a standard discretisation of a Poisson equation in two dimensions on a rectangular, uniform grid, the CPU-time for CG with RRB-preconditioning is not more than $O(N^{9/8})$. In [1] it is shown that an optimal order of convergence can be obtained by using such an incomplete decomposition and certain combinations of V-cycles and W-cycles. In [11] a different lumping procedure is used (e.g. a nine-point stencil is reduced to a five-point stencil using an approximation based on a linear interpolation) and the incomplete decomposition is combined with a multigrid method.

We have followed a different approach: after a renumbering of the unknowns according to the partitioning based on multigrid, $L$ and $U$ are obtained from a modified incomplete LU-decomposition based on a drop tolerance. In the sequel of this paper, this technique is referred to as Nested Grids ILU-decomposition (NGILU). This decomposition includes the modification of the main diagonal as described in [8]. The renumbering is similar to that of ILUM-factorization described in [12] in which the following idea is applied recursively: form the reduced system and apply a dropping strategy to this system. However, in the approach we follow, the renumbering is applied only once on the whole system of linear equations, and not each time when a reduced system has been formed.

In Section 2 the partitioning of the unknowns and the construction of the factors $L$ and $U$ is described. Some numerical results are given in Section 3. In Section 4, some conclusions are drawn, and some suggestions for future research are made.
2. The preconditioning technique

First we describe the renumbering of the unknowns. Consider a sequence of nested grids \( \Omega_1, \Omega_2, \ldots, \Omega_m \), where \( \Omega_\gamma \subseteq \Omega_{\gamma-1} \subseteq \cdots \subseteq \Omega_1 \). If all grids are uniform, \( \Omega_m \) has mesh size \( 2^{m-1} h \), where \( h \) is the mesh size of the finest grid \( \Omega_1 \). The set of unknowns at the \( m \)th level is now defined by \( \Omega_m = \Omega_m \setminus \Omega_{m+1} \), where \( \Omega_{m+1} = \emptyset \). If the numbering within the levels is lexicographical, and if Dirichlet boundary conditions are used, we obtain for the inner grid points of a rectangular 8 × 8-grid with constant mesh size:

\[
\begin{array}{cccccccc}
1 & 2 & 3 & 4 & 5 & 6 \\
7 & 28 & 8 & 29 & 9 & 30 \\
10 & 11 & 12 & 13 & 14 & 15 \\
16 & 31 & 17 & 36 & 18 & 32 \\
19 & 20 & 21 & 22 & 23 & 24 \\
25 & 33 & 26 & 34 & 27 & 35 \\
\end{array}
\]

The points with numbers 1 to 27 belong to the first level \( \Omega_1 \). Similarly, the two sets of points 28 to 35 and 36 belong to \( \Omega_2 \) and \( \Omega_3 \), respectively (note that the number of inner grid points in one direction does not necessarily have to be a power of 2). In [15] an algorithm for the generation of such a numbering is given in the more general case where the mesh is not uniform. This algorithm is based on a sequence of uniform grids which are laid over the physical domain determining to which level the points belong. The idea is similar to that of adaptive multigrid [9]. Numbering the unknowns as described above results in a system of linear equations which can be written as

\[
\begin{bmatrix}
A_{11} & A_{12} \\
A_{21} & A_{22}
\end{bmatrix}
\begin{bmatrix}
x_1 \\
x_2
\end{bmatrix} =
\begin{bmatrix}
b_1 \\
b_2
\end{bmatrix},
\]

where \( x_1 \) is the vector containing the unknowns of the first level \( \Omega_1 \), and \( x_2 \) those of the second grid \( \Omega_2 \). The partitioning of the matrix can be repeated for the matrix in the lower-right corner until we arrive at the coarsest grid. In order to obtain factors \( L \) and \( U \) which enable us to make an efficient implementation of the statement \( y := (LU)^{-1} z \) on supercomputers, it is advantageous to use an appropriate choice for the ordering of unknowns within each level. For example, one can use a red–black ordering.

2.1. Construction of the factors \( L \) and \( U \)

The preconditioning technique consists now of making a splitting \( A = LU + R \). In [14], several methods for choosing such a splitting were compared to each other, and it appeared that strategies based on a threshold parameter led to the best results. Therefore, we construct the splitting in such a way that the elements of the residual matrix \( r_{ij} \) all satisfy \( |r_{ij}| \ll e_{ij} \). This is realized by making an incomplete LU-decomposition of \( A \) based on the threshold parameter \( e_{ij} \) in a similar way as in [13]. First the main diagonal of \( A \) is scaled to unity. Next, the incomplete decomposition is
constructed row by row. Given the first \( i - 1 \) rows of \( L \) and \( U \), we construct row \( i \) of \( L \) and \( U \) from \( R = A - LU \) as

\[
 r_{ik} = a_{ik} - \min_{j=1}^{i} l_{ij} u_{jk}. \tag{1} \]

Suppose \( l_{ij} \) has been calculated for \( j < k \). If \( k < i \) one obtains from (1)

\[
 r_{ik} + l_{ik} u_{kk} = a_{ik} - \sum_{j=1}^{k-1} l_{ij} u_{jk}, \quad k < i. \tag{2} \]

If the absolute value of the right-hand side of this equation is less than \( \varepsilon_{ij} \), fill-in on position \( (i, k) \) is added on the main diagonal (Gustafsson’s modification [8]: by this choice a constant vector is in the kernel of \( R \)); otherwise \( l_{ik} \) is calculated from (2) together with \( r_{ik} = 0 \). With \( l_{ii} = 1 \), \( u_{ii} \) can be calculated from

\[
 u_{ii} = a_{ii} - \sum_{j=1}^{i-1} l_{ij} u_{ji}. \tag{3} \]

If \( k > i \), \( u_{ik} \) can be calculated in a similar way. In the special case of a positive definite, symmetric matrix \( A \), we demand that \( U = L^T \), and we can construct an incomplete Choleski-decomposition in a similar way.

2.2. Choice of the drop tolerance

In the special case where all mesh sizes are approximately the same, \( \varepsilon_{ij} \) is chosen as follows:

\[
 \varepsilon_{ij} = c^{m-1} \varepsilon^{(1)}, \quad \text{where} \quad \max(i,j) \in W_m. \tag{4} \]

Starting with \( \varepsilon^{(1)} \) in the first diagonal block, corresponding with level \( W_1 \), we let the drop tolerance decrease by multiplying with a positive factor \( c < 1 \) at each new level, as shown in Fig. 1.

This particular choice of the drop tolerance causes an increase of the fill-in per level, which is essential for the quality of the incomplete decomposition. For most problems, the choice of the parameters \( \varepsilon^{(1)} \) and \( c \) is not very critical. In 2D-problems \( c = 0.2 \) is a reasonable choice, but in 3D-problems the optimal value for \( c \) is smaller. This choice will be motivated by the analysis below.

\[
 \varepsilon^{(1)} \]

\[
 \varepsilon^{(1)}_c \]

Fig. 1. The drop tolerance in case of a (nearly) uniform grid.
Though the approach is heuristic, more insight in the behaviour of the algorithm is obtained in the course of this derivation. Assume that the system of linear equations \( Ax = b \) arises from some discretisation of a Poisson equation on a rectangular grid, in such a way that \( A \) is symmetric positive definite.

The criterion for choosing the drop tolerance is based on the following lemma.

Lemma 1. Let the incomplete Choleski-decomposition of the nonsingular matrix \( A \) be given by \( LL^T = A - R \). If both \( R \) and \( A - \alpha R \) are positive semidefinite for \( \alpha > 1 \), then the condition number of \( L^{-1}AL^{-T} \) is bounded by \( \frac{\alpha}{\alpha - 1} \).

Remark. Suppose that \( A \) is an \( M \)-matrix, and \( L \) results from a modified incomplete decomposition of \( A \) as described in [8]. In that case \( R \) has nonpositive elements outside the main diagonal and row sums zero. Hence from Gerschgorin’s theorem it follows that \( R \) is positive semidefinite.

Proof of Lemma 1. From \( LL^T = A - R \) it follows that \( L^{-1}AL^{-T} = I + L^{-1}RL^{-T} \), and since \( R \) is positive semidefinite, it follows that none of the eigenvalues of \( L^{-1}AL^{-T} \) is less than 1.

From \( A - \alpha R \) being positive semidefinite and \( R = A - LL^T \), it follows that \((1 - \alpha)L^{-1}AL^{-T} + \alpha I\) is positive semidefinite. This completes the proof of Lemma 1. \( \Box \)

In the following, we try to find a condition which assures that \( A - \alpha R \) is positive semidefinite for some \( \alpha > 1 \). Our starting point will be the necessary condition for this requirement that

\[ u_k^T(A - \alpha R)u_k \geq 0, \]

where \( u_k \) is a normalized eigenvector of \( A \). This condition is sufficient if the eigenvectors \( u_k \) are equal or sufficiently close to the eigenvectors of \( A - \alpha R \). As we need indices below for other purposes, we denote \( u_k \) by \( q \), and rewrite our condition in the form

\[ \alpha q^T R q \leq \lambda, \]  \( (5) \)

where \( \lambda \) is the eigenvalue of \( A \) corresponding to \( q \). Note that (5) is a strong requirement for low-frequency components. For example, for a nearly constant eigenvector with a very small eigenvalue \( \lambda \), we see that Gustafsson’s modification is necessary. Numerical experiments showed that \( R \) is of near block-diagonal form. Therefore, we restrict our analysis to the strict block-diagonal form

\[ R = \begin{bmatrix} R_{11} & \cdots & \cdots \cr \cdots & R_{22} & \cdots \cr \cdots & \cdots & \cdots \cr \cdots & \cdots & \cdots \cr \cdots & \cdots & \cdots \cr \cdots & \cdots & \cdots \cr \cdots & \cdots & \cdots \cr R_{rr} \end{bmatrix}. \]  \( (6) \)

Consider \( q^T R q \). From (6) it follows that

\[ q^T R q = \sum_{i=1}^{\gamma} q_i^T R_{ii} q_i, \]
where, using the same partitioning, \( q = [q_1, q_2, \ldots, q_r]^T \). Next suppose that
\[
q_i^T R_{ii} q_i \leq \mu_i q_i^T q_i,
\]
then
\[
q^T R q = \sum_{i=1}^{r} q_i^T R_{ii} q_i \leq \sum_{i=1}^{r} \mu_i q_i^T q_i.
\]

Now, for all \( i \) it holds that \( q_i \) is a restriction of \( q \) to a coarser grid. For a smooth eigenvector \( q \), one can find a smooth eigenfunction \( u(x, y) \) of the partial differential equation, such that the components of the vector \( q \) are given by the value of \( u(x, y) \) in the grid points. In that case, both \( q^T q \) and \( q_i^T q_i \) represent, up to a factor, an approximation to the integral
\[
\int_0^1 \int_0^1 u(x, y) \, dx \, dy.
\]
The ratio between these two approximations can be found by substituting a constant in the respective integration rules, thus it is approximately equal to the ratio of the lengths of the vectors \( q \) and \( q_i \). Hence \( q_i^T q_i \approx 2^{-i} \) (e.g. \( q_1 \) has about half the length of \( q \)). Furthermore, suppose that there exists a real \( \beta \) such that \( \mu_i = \mu_0 \beta^i \) and \( q_i^T q_i \leq 2^{-i} \), then
\[
q^T R q \leq \mu_0 \sum_{i=1}^{r} 2^{-i} \beta^i = \frac{1}{2} \mu_0 \beta \sum_{i=1}^{r} (\frac{1}{2} \beta)^{i-1} = \frac{1}{2} \mu_0 \beta \frac{1 - (\frac{1}{2} \beta)^r}{1 - \frac{1}{2} \beta}.
\]

Now if \( \beta < 2 \) then for all \( r \)
\[
q^T R q \leq \frac{\mu_0 \beta}{2 - \beta}.
\]
Hence, under the assumptions made, condition (5) is satisfied if
\[
\frac{\mu_0 \beta}{2 - \beta} < \frac{\lambda}{\alpha}, \quad \alpha > 1.
\]

It should be noted that the condition \( \beta < 2 \) allows a growth of \( \mu_i \) with \( i \). Furthermore, numerical experiments showed that far enough away from the boundaries the blocks \( R_{ii} \) have a very regular structure and constant coefficients. This allows us to use Fourier analysis in order to obtain \( \mu_i \). The use of Fourier analysis for estimating the performance of preconditioners is, for example, described in [5] and [6]. Note also that for our model problem the eigenvectors \( q \) are in fact Fourier components with a special choice for the frequency.

Hence, we approximate \( R_{ii} \) by a linear combination of elementary symmetric finite-difference stencils. In our current algorithm all contributions to \( R_{ii} \) are balanced. More precise, define \( \mu_i = \sum \delta_j \) where the sum is over all stencils, then the eigenvalues \( \delta_j \) satisfy the condition
\[
\delta_j < \delta_{\text{max}} = \delta \lambda \quad \text{for all } j.
\]

Hence the stencil is allowed in \( R_{ii} \) if (7) holds, otherwise it should be included in the decomposition. In the numerical experiments we observed a moderate growth of fill-in with increasing \( i \). Hence \( \mu_i \)
will grow with the number of significant stencils (with $\delta_j$ close to $\delta_{\text{max}}$) but this increase is not so large that $\beta = \mu_{i+1}/\mu_i \geq 2$. In the following we will make condition (7) more specific.

Suppose that $A$ stems from a standard five-point discretisation of a Poisson equation on an $M_x \times M_y$-grid with mesh size $h$ and $k$ in the horizontal and vertical direction, respectively. When the diagonal is scaled to unity, $A$ has the five-point stencil

$$\begin{bmatrix}
\frac{1}{2(h^2 + k^2)} & -h^2 & -h^2 \\
-k^2 & 2(h^2 + k^2) & -k^2 \\
-h^2 & -k^2 & \frac{1}{2(h^2 + k^2)}
\end{bmatrix}.$$

The eigenvalues for a Fourier component $\exp(i(f_1 x + f_2 y))$ of this stencil are

$$\lambda_A(f_1, f_2) = \frac{1}{2(h^2 + k^2)} [2(h^2 + k^2) - 2k^2 \cos f_1 h - 2h^2 \cos f_2 k]$$

$$= \frac{2}{(h^2 + k^2)} [k^2 \sin^2(\frac{1}{2} f_1 h) + h^2 \sin^2(\frac{1}{2} f_2 k)]. \quad (8)$$

Here $|f_1 h|, |f_2 k| < \pi$. Suppose that $R_{ii}$ is the sum of stencils of the form

$$\varepsilon \begin{bmatrix}
-1 & -1 \\
4 & 0 \\
-1 & -1
\end{bmatrix},$$

with mesh sizes $nh$ and $mk$ in the horizontal and vertical direction, respectively. Here, $n$ or $m$ may be zero! Note that we may take the sum of the coefficients equal to zero, because the row sums of $R$ are zero. The eigenvalue $\delta_j$ of this stencil is given by

$$\delta_j(f_1, f_2) = 4\varepsilon [1 - \cos f_1 nh \cos f_2 mk]$$

$$= 8\varepsilon \left[ \sin^2(\frac{1}{2} n f_1 h) + \sin^2(\frac{1}{2} m f_2 k) - 2\sin(\frac{1}{2} n f_1 h) + \sin^2(\frac{1}{2} m f_2 k) \right]. \quad (9)$$

Now $\delta_j(f_1, f_2)$ can be bounded from above by

$$\delta_j(f_1, f_2) \leq 8\varepsilon \left[ \sin^2(\frac{1}{2} n f_1 h) + \sin^2(\frac{1}{2} m f_2 k) \right]$$

$$\leq 8\varepsilon [n^2 \sin^2(\frac{1}{2} f_1 h) + m^2 \sin^2(\frac{1}{2} f_2 k)]. \quad (10)$$

Note that in the first step equality occurs only if $n f_1 h$ or $m f_2 k$ is a multiple of $\pi$, and in the second step if $f_1 = f_2 = 0$. The second step holds only because of the limitations put on $f_1 h, f_2 k$ given below Eq. (8).
Condition (7) is, in this specific case, \( \delta_j(f_1, f_2) < \delta\lambda_A(f_1, f_2) \). This is certainly true if it holds for \( \delta_j \) replaced by the upper bound (10), hence if

\[
\varepsilon < \frac{\delta(k^2 \sin^2(\frac{1}{2} f_1 h) + h^2 \sin^2(\frac{1}{2} f_2 k))}{4(h^2 + k^2)(n^2 \sin^2(\frac{1}{2} f_1 h) + m^2 \sin^2(\frac{1}{2} f_2 k))}.
\]

In its turn this condition is certainly fulfilled if

\[
\varepsilon < \frac{\delta h^2 k^2}{4(h^2 + k^2) \min \left( \frac{1}{(nh)^2}, \frac{1}{(mk)^2} \right)}.
\]

We obtained slightly better results with the following choice:

\[
\varepsilon < \frac{\delta h^2 k^2}{4(h^2 + k^2)(n^2 h^2 + m^2 k^2)}.
\]

Note that in the special case where the grid is uniform, \( h = k \) and \( n = m \) and we obtain a drop tolerance as shown in Fig. 1 with \( c = \frac{1}{4} \). In our implementation with varying meshes we use

\[
\varepsilon_{ij} = \varepsilon \frac{h_{ij}^2 k_{ij}^2}{(h_{ij}^2 + k_{ij}^2) \rho_{ij}^2}.
\]

Herein \( h_{ij} \) and \( k_{ij} \) are the minimum of the mesh sizes in, respectively, horizontal and vertical direction at points \( i \) and \( j \), \( \rho_{ij} \) is the distance between the two grid points with numbers \( i \) and \( j \), and \( \varepsilon \) is a parameter which has to be chosen in advance.

3. Numerical experiments

In this section, we demonstrate NGILU-preconditioning by showing the results of three examples. When a lexicographical ordering of the unknowns is used, and the sparsity pattern of \( L + U \) is the same as that of \( A \), this is indicated as standard (M)ILU. When the sparsity pattern of the factors \( L \) and \( U \) is based on a threshold parameter \( \varepsilon \), this is indicated as (M)ILU(\( \varepsilon \)). As stopping criterion we used

\[
\| L^{-1}(b - AU^{-1}\tilde{x}^{(m)}) \|_2 < 10^{-6}\| L^{-1}(b - AU^{-1}\tilde{x}^{(0)}) \|_2,
\]

where \( \tilde{x}^{(m)} = Ux^{(m)} \).

Example 1. The first example shows the results of a Nested Grids Incomplete Choleski-decomposition (NGIC) applied to a matrix which can be represented by a nine-point stencil. Consider the Poisson equation on the unit square \([0, 1] \times [0, 1]\) with condition \( \frac{\partial u}{\partial n} = 0 \) on every boundary. This problem is of interest for unsteady incompressible Navier–Stokes solvers, where at each time step the pressure has to be computed. The Poisson equation is discretised over a rectangular \((M - 1) \times (M - 1)\)-grid with constant mesh size, using a standard five-point stencil. Within the
separate levels we use a red–black ordering. Hence after renumbering and scaling we obtain the following system of linear equations

\[
\begin{pmatrix}
I & C \\
C^T & I
\end{pmatrix}
\begin{pmatrix}
x_1 \\
x_2
\end{pmatrix} =
\begin{pmatrix}
b_1 \\
b_2
\end{pmatrix},
\]

which can be solved by first solving the reduced system \((I - C^TC)x_2 = b_2 - C^Tb_1\) and then calculating \(x_1 = b_1 - Cx_2\). The Schur-complement \((I - C^TC)\) can be represented by a nine-point stencil (see, e.g., [7]). Table 1 shows the results of NGIC-decomposition combined with the Conjugate Gradient method (NGICCG) applied to the reduced system. For the drop tolerance we used (4) with \(\epsilon^{(1)} = c = 0.2\). The second row shows that the number of CG-iterations does not increase with mesh-refinement. The third row shows that the number of non-zero elements in \(L\) divided by \(M^2\) hardly increases. As a consequence, the number of floating point operations is almost a constant times the number of degrees of freedom.

Example 2. This test problem is a simplified aquifer problem which is taken from [18]. The nonsymmetric system of linear equations comes about after the discretisation of the steady convection–diffusion equation

\[
-\frac{\partial}{\partial x} \left( A \frac{\partial}{\partial x} u \right) - \frac{\partial}{\partial y} \left( A \frac{\partial}{\partial y} u \right) + B \frac{\partial}{\partial x} u = F
\]

on the square \([0, 1] \times [0, 1]\). The function \(B(x, y)\) is equal to \(2e^{2(x^2 + y^2)}\), and the diffusion coefficient \(A(x, y)\) is strongly discontinuous as shown in Fig. 2. The dashed area indicates the region in which \(A(x, y) = 10000\). The right-hand side \(F(x, y)\) is zero everywhere, except for the small square in the center, where \(F(x, y) = 100\).

As a preconditioning technique we used an NGILU-decomposition with the drop tolerance as in (4) with \(\epsilon^{(1)} = c = 0.2\). Fig. 3 shows that the convergence rate of Bi-CGSTAB is excellent, even when the coefficients in the PDE are strongly discontinuous. From the results of Fig. 4 it appears that the convergence behaviour is relatively smooth, which is advantageous to the construction of stopping criteria when the linear solver is used as an inner-iteration method, for example, within some Newton method.

Example 3. The convergence rate of multigrid algorithms based on point relaxation smoothers deteriorates for problems with strong anisotropies (see, e.g., [3]). Anisotropic discrete operators arise, for example, in problems in which the differential operator is discretised on highly stretched

<table>
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<th>(M)</th>
<th>65</th>
<th>129</th>
<th>257</th>
<th>513</th>
<th>1025</th>
</tr>
</thead>
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<tr>
<td># iterations</td>
<td>7</td>
<td>7</td>
<td>7</td>
<td>7</td>
<td>7</td>
</tr>
<tr>
<td># nonz./(M^2) for it.</td>
<td>3.6</td>
<td>4.0</td>
<td>4.3</td>
<td>4.5</td>
<td>4.6</td>
</tr>
<tr>
<td># flops/(M^2) for it.</td>
<td>204</td>
<td>217</td>
<td>226</td>
<td>233</td>
<td>236</td>
</tr>
</tbody>
</table>
grids. Therefore, we tested the convergence of NGICCG for the system of linear equations which arises after the discretisation of a Poisson equation on a highly stretched rectangular grid. We took the same test problem as in [3]. We consider \(-\Delta u(x, y) = f(x, y)\) on the unit square with the boundary condition \(\partial u/\partial n = 0\) along \(x = 0\) and \(y = 0\), and Dirichlet boundary conditions along
Table 2
Numerical results for Example 3

<table>
<thead>
<tr>
<th>$h_{\max}/h_{\min}$</th>
<th>$M = 32$</th>
<th>$M = 64$</th>
<th>$M = 128$</th>
<th>$M = 256$</th>
</tr>
</thead>
<tbody>
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<td>6 (9.7)</td>
<td>7 (9.7)</td>
<td>7 (9.8)</td>
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<tr>
<td>100</td>
<td>5 (10.8)</td>
<td>5 (11.5)</td>
<td>7 (11.2)</td>
<td>9 (11.2)</td>
</tr>
<tr>
<td>$10^3$</td>
<td>4 (12.1)</td>
<td>5 (16.3)</td>
<td>6 (15.7)</td>
<td>10 (15.1)</td>
</tr>
<tr>
<td>$10^4$</td>
<td>3 (11.6)</td>
<td>4 (15.6)</td>
<td>5 (22.0)</td>
<td>9 (25.1)</td>
</tr>
<tr>
<td>$10^5$</td>
<td>3 (11.0)</td>
<td>3 (15.3)</td>
<td>4 (21.7)</td>
<td>7 (28.4)</td>
</tr>
<tr>
<td>$10^6$</td>
<td>3 (10.8)</td>
<td>3 (15.0)</td>
<td>5 (22.4)</td>
<td>6 (31.0)</td>
</tr>
</tbody>
</table>

the other boundaries. An exponential stretching of an $(M + 1) \times (M + 1)$-grid was used in both coordinate directions. The stretching was done in such a way that the minimum mesh sizes occur near the boundaries with Neumann boundary conditions, and the largest mesh sizes occur near the other boundaries. Table 2 lists the number of iteration steps of CG for various choices of $M$ and $h_{\max}/h_{\min}$. For the drop tolerance $\varepsilon_{ij}$ we used (11) with $\varepsilon = 0.2$. The average number of entries in one row of $L$ is given between brackets. When a strong refinement of the grid is used, the factor $L$ contains more elements, but the number of CG iteration steps combined with the resulting preconditioner remains very small.

4. Conclusions

In this paper, we have described a preconditioning technique (NGILU) which uses a partitioning of the unknowns based on a similar sequence of grids as in multigrid. Renumbering the unknowns according to this partitioning enables us to construct an incomplete LU-decomposition which can be used in eliminating effectively both high- and low-frequency errors. The factors $L$ and $U$ are obtained from an incomplete decomposition based on a drop tolerance. Results are presented from the above method applied to test problems including strongly varying coefficients, advective terms and grid refinement. In all cases, the method is much cheaper than standard (M)ILU. This difference is more pronounced for the really difficult problems and increases with the dimension. From numerical experiments we observed that, for many other problems, an NGILU-decomposition is superior to a standard (M)ILU-preconditioning (see [15] and [16]).

In the present method, the renumbering of the unknowns is based on a rectangular grid. The authors are currently working on an algorithm which produces a renumbering using only the matrix entries. This variant enables us to use the NGILU concept over a wide range of problems: for example, for systems of linear equations that arise after discretisation of PDEs on irregular domains using some finite element package. The first results of this method look very promising.

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References


